

## UNIT-1: Introduction & Analysis of Algorithms

### Algorithm

An **algorithm** is a finite sequence of well-defined steps used to solve a problem.

### Characteristics of an Algorithm

1. **Input** – Takes zero or more inputs.
  2. **Output** – Produces at least one output.
  3. **Definiteness** – Each step is clear and unambiguous.
  4. **Finiteness** – Terminates after a finite number of steps.
  5. **Effectiveness** – Each step is basic and feasible.
  6. **Generality** – Applicable to a class of problems.
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### Analysis of Algorithm

Algorithm analysis determines **efficiency** in terms of **time and space**.

### Asymptotic Analysis

Used to analyze algorithm behavior for **large input sizes (n)**.

### *Complexity Bounds*

1. **Big-O (O)** – Upper bound (Worst case)
2. **Big-Ω (Ω)** – Lower bound (Best case)
3. **Big-Θ (Θ)** – Tight bound (Average case)

Example:

Linear Search

- Best case:  $\Omega(1)$
  - Worst case:  $O(n)$
  - Average case:  $\Theta(n)$
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## Best, Average, and Worst Case Analysis

- **Best Case:** Minimum time taken
  - **Average Case:** Expected time for random input
  - **Worst Case:** Maximum time taken
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## Performance Measurement

1. **Time Complexity** – Number of operations
  2. **Space Complexity** – Memory used
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## Time–Space Trade-off

An algorithm may use:

- More memory to reduce time (e.g., hashing)
  - More time to save memory (e.g., brute force)
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## Analysis of Recursive Algorithms

Uses **recurrence relations**.

### Substitution Method

- Guess the solution
- Prove by induction

### Recursion Tree Method

- Represent recursive calls as a tree
- Sum cost at each level

### Master's Theorem

Used for recurrences of form:

$$T(n) = aT(n/b) + f(n)$$

Three cases:

- 1.
- 2.
- 3.

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## UNIT-2: Divide and Conquer & Heaps

### Divide and Conquer Paradigm

Steps:

1. Divide problem into sub-problems
2. Conquer recursively
3. Combine results

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### Binary Search

- Search in sorted array
- Time:  $O(\log n)$

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### Merge Sort

- Divide array, sort recursively, merge
- Time:  $O(n \log n)$
- Stable sorting

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### Quick Sort

- Choose pivot
- Partition array
- Average:  $O(n \log n)$
- Worst:  $O(n^2)$

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## Linear Time Selection

Find k-th smallest element using **Median of Medians**

- Time:  $O(n)$

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## Strassen's Matrix Multiplication

- Reduces multiplication operations
- Time:  $O(n^{2.81})$

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## Karatsuba Algorithm

- Fast multiplication of large numbers
- Time:  $O(n^{1.585})$

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## Heap

A **complete binary tree** satisfying heap property.

### Min Heap

- Parent  $\leq$  children

### Max Heap

- Parent  $\geq$  children

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## Build Heap

- Convert array into heap
- Time:  $O(n)$

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## Heap Sort

1. Build heap
  2. Repeatedly remove root
- Time:  $O(n \log n)$
  - In-place sorting

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## UNIT–3: Algorithm Design Techniques

### Brute Force

- Try all possibilities
- Simple but inefficient
- Example: Linear search

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### Greedy Algorithm

- Make locally optimal choice

### Greedy Examples

1. **Minimum Cost Spanning Tree**
    - Prim's Algorithm
    - Kruskal's Algorithm
  2. **Knapsack (Fractional)**
    - Choose highest profit/weight ratio
  3. **Job Sequencing**
    - Maximize profit with deadlines
  4. **Huffman Coding**
    - Optimal prefix code for compression
  5. **Single Source Shortest Path**
    - Dijkstra's algorithm
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## Backtracking

- Try possible solutions
  - Undo when constraint violated
  - Example: N-Queens
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## Branch and Bound

- Optimization version of backtracking
  - Uses bounds to prune search space
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# UNIT-4: Dynamic Programming & Heuristics

## Dynamic Programming (DP)

Used when:

- Overlapping sub-problems
  - Optimal substructure
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## DP vs Divide and Conquer

DP	Divide & Conquer
Stores results	No storage
Avoids recomputation	Recomputes
Bottom-up	Top-down

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## DP Applications

1. **Fibonacci Series**
2. **Matrix Chain Multiplication**
3. **0-1 Knapsack**

4. **Longest Common Subsequence (LCS)**
  5. **Travelling Salesman Problem**
  6. **Rod Cutting**
  7. **Bin Packing**
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## Heuristics

- Approximate solutions
- Faster than exact algorithms
- Used in NP-hard problems

## Characteristics

- Not always optimal
- Problem-specific
- Efficient

## Application Domains

- Scheduling
  - Routing
  - AI problems
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# UNIT-5: Graph & Tree Algorithms

## Graph Representation

1. **Adjacency Matrix**
  2. **Adjacency List**
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## Traversal Algorithms

### DFS

- Uses stack/recursion
- Depth-wise traversal

## BFS

- Uses queue
  - Level-wise traversal
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## Shortest Path Algorithms

### Bellman-Ford

- Handles negative weights
- Time:  $O(VE)$

### Dijkstra's Algorithm

- Greedy approach
- Uses priority queue
- Time:  $O(E \log V)$

### Floyd-Warshall

- All-pairs shortest path
  - Time:  $O(n^3)$
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## Other Graph Algorithms

- **Transitive Closure**
  - **Topological Sorting**
  - **Network Flow (Ford-Fulkerson)**
  - **Connected Components**
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## UNIT-6: NP Problems & Advanced Algorithms

### Tractable vs Intractable Problems

- **Tractable:** Polynomial time



- **Intractable:** Exponential time
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## Computability

Determines whether a problem is solvable by algorithm.

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## Complexity Classes

- **P:** Solvable in polynomial time
  - **NP:** Verifiable in polynomial time
  - **NP-Complete:** Hardest problems in NP
  - **NP-Hard:** At least as hard as NP-Complete
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## Cook's Theorem

- SAT is NP-Complete
  - Foundation of NP-Completeness
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## Standard NP-Complete Problems

- Travelling Salesman
  - Knapsack
  - Vertex Cover
  - Clique
  - Hamiltonian Cycle
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## Reduction Techniques

Transform one problem to another in polynomial time.

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## Approximation Algorithms

- Near-optimal solutions
  - Polynomial time
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## Randomized Algorithms

- Use random numbers
- Faster on average
- Example: Randomized Quick Sort